

# Bypassing Combinatorial Protections

Polynomial-Time Algorithms for Single-Peaked Electorates

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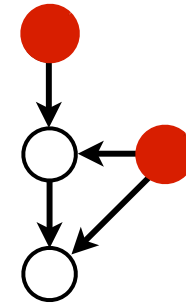
PREFERENCE AGGREGATION IN MULTIAGENT SYSTEMS

# Voting Rules

- $C = \{a, b, c, \dots\}$  is a finite set of **candidates or alternatives**
- A **voting rule**  $f$  maps a vector  $V = (v_1, v_2, \dots, v_n)$  of **votes** to a non-empty **subset**  $f(V) \subseteq C$  of candidates
  - ▶ ignores tie-breaking
- Ranking-based voting rules
  - ▶ each vote is a complete **ranking** of the candidates:  $v_i = [b \succ_i a \succ_i c]$
- Approval-based voting rules
  - ▶ each vote is a **set** of “approved” candidates:  $v_i = \{a, b\}$

# Voting rules (2)

- Ranking-based voting rules
  - ▶ Many rules are defined via pairwise comparisons (**majority graphs**)
  - ▶ **Weak Condorcet winners**: all candidates without pairwise defeats
  - ▶ A **weakCondorcet rule** is a rule that precisely returns all weak Condorcet winners whenever at least one exists
  - ▶ **Lull's rule** yields all candidates with minimal number of pairwise defeats
  - ▶ Young, Kemeny, Dodgson, Maximin, Fishburn, Schwartz, ...
- **Approval** voting
  - ▶ yields all candidates with maximal number of approvals



# Swaying Elections

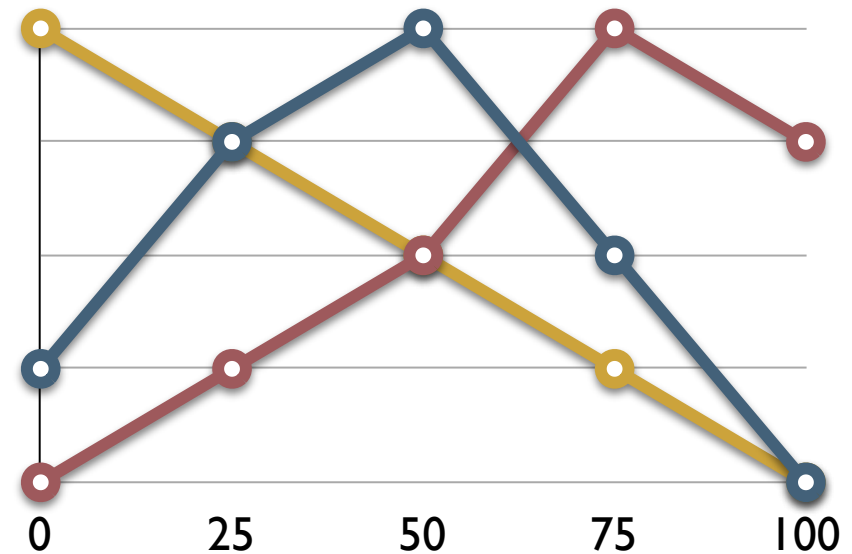
- People may try to **influence the outcome** of an election by
  - ▶ manipulating the voters' preferences (e.g., bribery, campaigning, strategic manipulation)
  - ▶ changing the election's structure (e.g., introducing primaries, adding/deleting voters and/or candidates)
- All voting rules are **vulnerable** to at least some of these attacks
  - ▶ But: Attacker's task may be **computationally intractable** due to **combinatorial** challenges (e.g., covering or partition problems) [BTT 1989]
- Are such computational protections meaningful **in practice**?
  - ▶ NP-hardness is a worst-case measure
  - ▶ heuristics that find successful manipulations in “most” instances
  - ▶ approximation algorithms



# Single-Peaked Preferences

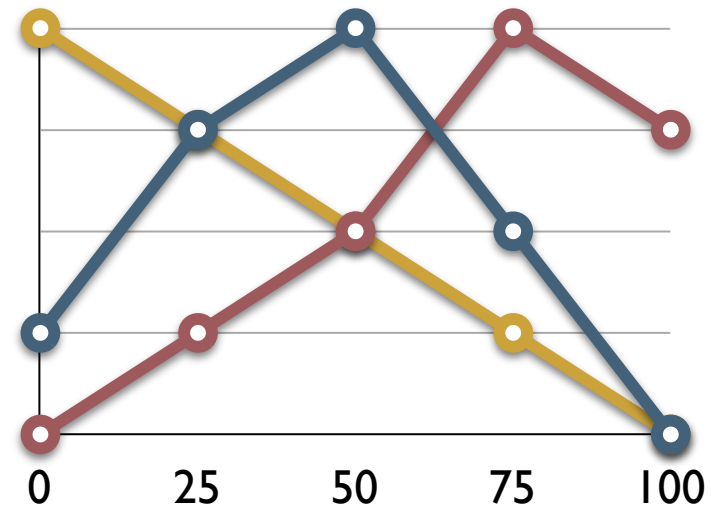
- What should be the **registration fee** for COMSOC 2010?
  - ▶ Candidates: €0, €25, €50, €75, €100

⊙ Jörg	⊙ Vince	⊙ Markus
€50	€75	€0
€25	€100	€25
€75	€50	€50
€0	€25	€75
€100	€0	€100



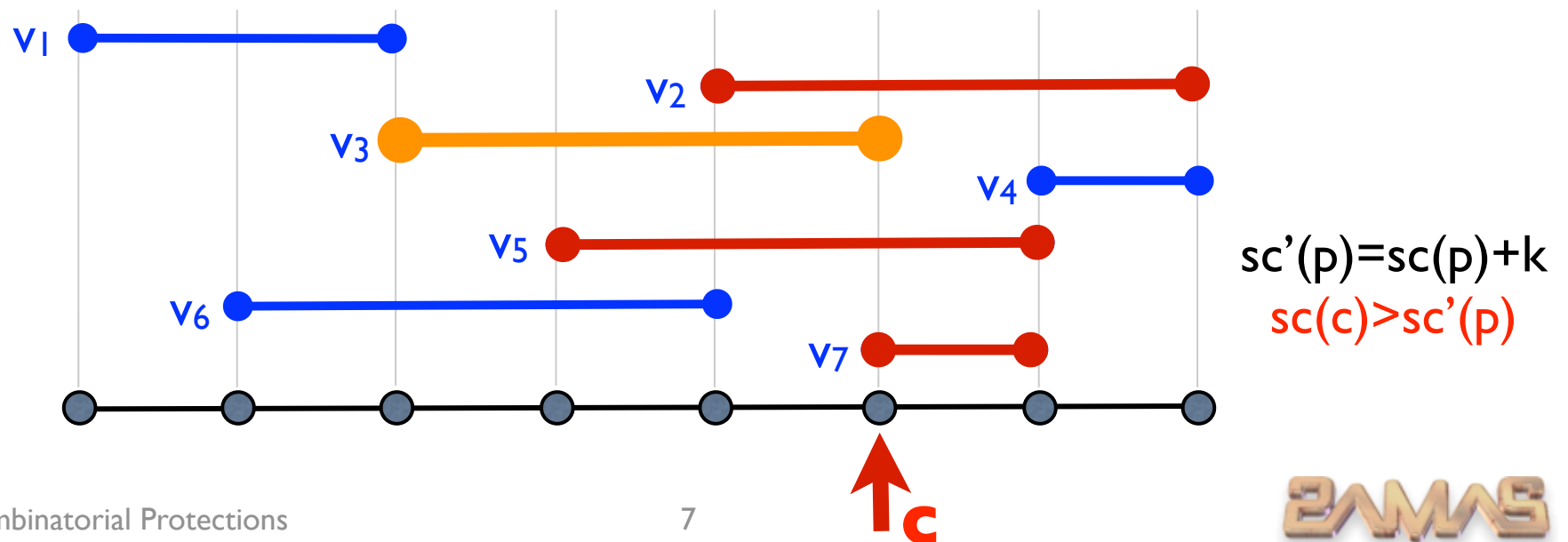
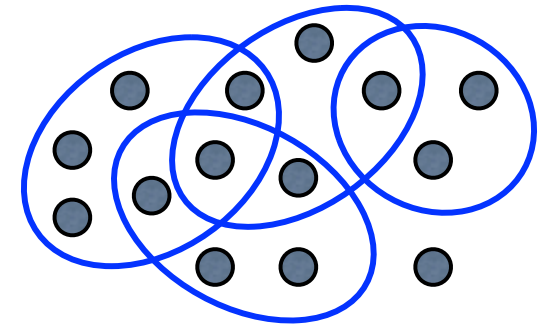
# Single-Peaked Preferences (2)

- Preferences are **single-peaked** iff there exists a linear ordering over  $C$  such that **if  $b$  lies between  $a$  and  $c$** , then  $(a \succ_i b \Rightarrow b \succ_i c)$  for all voters  $i$ 
  - ▶ natural variant for approval votes: approved candidates form an **interval**
  - ▶ Popular model in political science
    - **left-right** political spectrum
  - ▶ Single-peakedness can be **checked** in polynomial time [BT 1986]
  - ▶ **weak Condorcet** winners always exist (nice characterization in terms of **median voters**)



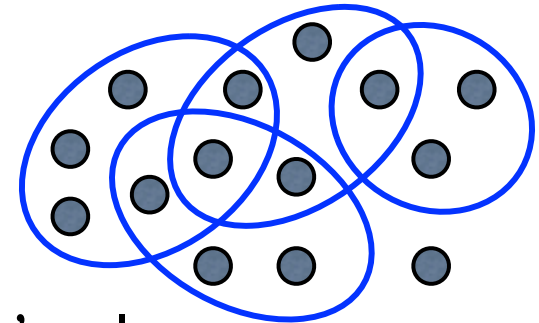
# Bribery

- Is it possible to **bribe at most  $k$  voters** such that  $p$  wins?
  - ▶ NP-complete for approval voting
    - reduction from X3C [FHH 2009]
  - ▶ Theorem: For **single-peaked** electorates, **approval bribery** is in P



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  - ▶ NP-hard for Lull's rule and [FHH 2009] Kemeny's rule
  - ▶ Theorem: For **single-peaked** electorates, this problem is in P for **all weakCondorcet rules**
    - e.g., Fishburn, Maximin, Young, Lull, *Kemeny*, *Schwartz*, *Nanson*, etc.





# Control

- Is it possible to **add/delete  $k$  voters** such that  $p$  wins?
  - ▶ NP-hard for Kemeny's rule and Young's rule
  - ▶ Theorem: For **single-peaked electorates**, both problems are in P for **all weakCondorcet rules**
- Can the set of voters be **partitioned** into two subsets (primary elections) such that  $p$  wins the final election?
  - ▶ NP-complete for Lull's rule [FHHR 2009]
  - ▶ Theorem: For single-peaked electorates, this problem is in P for **all weakCondorcet rules**

2x	1x	2x	1x
b	a	c	a
a	b	a	c
c	c	b	b

# Manipulation

- Constructive coalition weighted manipulation problem:  
Is it possible to **set the preferences of manipulative voters** such that  $p$  wins?
- We **completely characterize all scoring rules** where CCWM is in P or NP-complete for single-peaked electorates, respectively

# Conclusion

- It has been shown in previous work that various **manipulative attacks** on voting rules are computationally intractable
- In many realistic settings preferences may be assumed to be **single-peaked**
- The preference profiles constructed in many hardness proofs are **so intricate that they cannot be realized** by single-peaked electorates
- Good news: **Young, Kemeny, and Dodgson winners** can be computed in P for single-peaked electorates