

The Complexity of Bribery in Elections

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Context

- **Manipulation** [BTT89, many others]: Can voter(s) make Hillary a winner by voting strategically?
- **Control** [BTT92, not many others]: Can the “chair” of an election make Hillary win by addition/suppression/partition of voters/candidates?
- **Bribery**: Can a briber make Hillary win by spending a certain amount of money to bribe voters?

Setting

1. E-bribery (E an election system):

- **Given:** A set of candidates C , a set of voters V specified by their preference lists, a preferred candidate p in C , and an integer k .
- **Question:** Is it possible to make p a winner by changing the preference lists of at most k voters?

2. E-\$bribery:

- Same except that each voter has a price and that k is now a spending limit.

3. E-weighted-bribery and E-weighted-\$bribery:

- Same as 1 and 2, except that voters have weights.

Plurality-\$bribery Example

Voters:

1. John > Pat > Hillary > Ralph (\$10)
2. John > Pat > Hillary > Ralph (\$10)
3. John > Hillary > Ralph > Pat (\$10)
4. John > Ralph > Pat > Hillary (\$10)
5. Ralph > Hillary > John > Pat (\$1)
6. Pat > John > Ralph > Hillary (\$2)
7. Hillary > Ralph > John > Pat (\$1)

Can we make Hillary a winner by spending at most \$11?

Plurality-\$bribery Example

Voters:

1. Hillary > Ralph > John > Pat (\$10)
(John > Pat > Hillary > Ralph)
2. John > Pat > Hillary > Ralph (\$10)
3. John > Hillary > Ralph > Pat (\$10)
4. John > Ralph > Pat > Hillary (\$10)
5. Hillary > John > Pat > Ralph (\$1)
(Ralph > Hillary > John > Pat)
6. Pat > John > Ralph > Hillary (\$2)
7. Hillary > Ralph > John > Pat (\$1)

And Hillary is a winner (for only \$11).

Results I: Plurality

| | unweighted | weighted |
|-----------|------------|-----------------|
| No prices | P (a) | P (b) |
| \$ | P (c) | NP-complete (d) |

(a) Easy greedy.

(b)/(c) Trickier.

(d) Even for just two candidates. But if weights or prices are unary, we drop down to P.

Election Systems: Scoring Protocols

Scoring Protocols (for m candidates) are described by an integer vector $\alpha = (\alpha_1, \dots, \alpha_m)$ ($\alpha_1 \geq \dots \geq \alpha_m \geq 0$); i 'th position gives α_i (times weight of voter) points. The candidate(s) with the most points win.

Particularly important scoring protocols:

- plurality: $\alpha = (1, 0, 0, \dots, 0)$
- veto: $\alpha = (1, 1, \dots, 1, 0)$
- k -approval: $\alpha = (1^k, 0^{m-k})$
- Borda count: $\alpha = (m-1, m-2, \dots, 0)$.

Results II: Dichotomy Theorems

- **Goal:** Completely characterize when bribery is hard and easy.
- **Theorem A:** For each scoring protocol $\alpha = (\alpha_1, \dots, \alpha_m)$, if $\alpha_1 = \alpha_m$, then α -weighted-bribery is in P; otherwise, it is NP-complete.
- **Theorem B:** For each scoring protocol $\alpha = (\alpha_1, \dots, \alpha_m)$, if $\alpha_2 = \alpha_m$, then α -weighted-bribery is in P; otherwise, it is NP-complete.
- **Method:** By linking the complexity of bribery to the complexity of (restricted versions of) manipulation using (a modified version of) the manipulation dichotomy result from [HH05] (see also PR 2006 and CS/CLS 2005 combined version).

Results II (continued): Dichotomy Theorems

Comments:

- Theorem B clearly implies that veto-weighted-bribery is NP-complete, even for 3 candidates. However, weights are crucial, since veto-bribery is in P.
- Theorem B was proven by connecting the complexity of manipulation and bribery for scoring protocols. However, this won't work in all settings:
 - approval-bribery is NP-complete, but approval-manipulation is in P.
 - We have designed an artificial voting system for which bribery is easy, but manipulation is NP-complete.

Summary and Conclusions

- We introduced and determined the complexity of bribery for various electoral systems (many more results and all the proofs are available in our 49-page UR TR).
- We related bribery to manipulation to obtain complexity transfer and meta-results.
- We've seen that hardness results are “fragile.”

Future Directions

- Other election systems.
- Approximation algorithms.
- Average case behavior.
- Incomplete information.
- More complicated bribe structures.
- Competing bribers.