

# Hybrid Voting Protocols and Hardness of Manipulation

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# Manipulation: Example

- 99 voters, 3 candidates (Red, Blue, Green).
  - 49 voters:  $R > B > G$ .
  - 48 voters:  $B > R > G$ .
  - 2 voters (Edith and Helger):  $G > B > R$ .
- Aggregation rule: Plurality
  - each voter casts a vote for one candidate.
  - the candidate with the largest number of votes wins.
  - draws are resolved by a coin toss.

# What Will Edith and Helger Do?

R: 49 votes

B: 48 votes



G > B > R



If I vote for G, R will get elected, so I'd rather vote for B

If Edith and Helger vote B > G > R, they can guarantee that B is elected

# Why Manipulation Is Bad

- Aggregation rules are designed with certain **social welfare** criteria in mind.
- Misrepresentation of preferences results in a **suboptimal** choice w.r.t. these criteria.
- Encourages dishonesty...

# What If We Change Aggregation Rule?

- Single Transferable Vote:



# Formal Setup

- $n$  voters
- $m$  candidates  $c_1, \dots, c_m$
- Preference of a voter  $i$ :  
a permutation  $\pi_i$  of  $c_1, \dots, c_m$   
(best to worst).
- Aggregation rule  $S$ :  
$$\pi_1, \dots, \pi_n \longrightarrow c_j.$$

# Voting Schemes: Examples

- Borda: a candidate gets
  - $m$  points for each voter who ranks him 1<sup>st</sup>,
  - $m-1$  point for each voter who ranks him 2<sup>nd</sup>, etc.
- Copeland:
  - candidate that wins the largest # of pairwise elections
- Maximin:
  - $c$ 's score against  $d$ : # of voters that prefer  $c$  to  $d$ ;
  - $c$ 's # of points: min score in any pairwise election.
- many, many others...

# Voting Schemes: Properties

- **Pareto-optimality**: if everyone prefers **a** to **b**, **b** does not win
- **Condorcet-consistency**: if there is a candidate that wins every **pairwise** election, this candidate wins
- **Majority**: if there is a candidate that is ranked first by a **majority** of voters, this candidate wins
- **Monotonicity**: it is impossible to cause a winning candidate to lose by moving it **up** in one's vote

**Arrow's** theorem: there is no perfect scheme



# Manipulation: Definition

- A voter  $i$  can manipulate a voting scheme  $S$  if there is

- a preference vector

$$\pi = (\pi_1, \dots, \pi_i, \dots, \pi_n)$$

- a permutation  $\pi_i'$  s.t.

$$S(\pi_1, \dots, \pi_i', \dots, \pi_n) >_i S(\pi).$$

**Theorem** (Gibbard-Satterthwaite, 1971):  
every non-dictatorial aggregation rule  
with  $\geq 3$  candidates is manipulable.

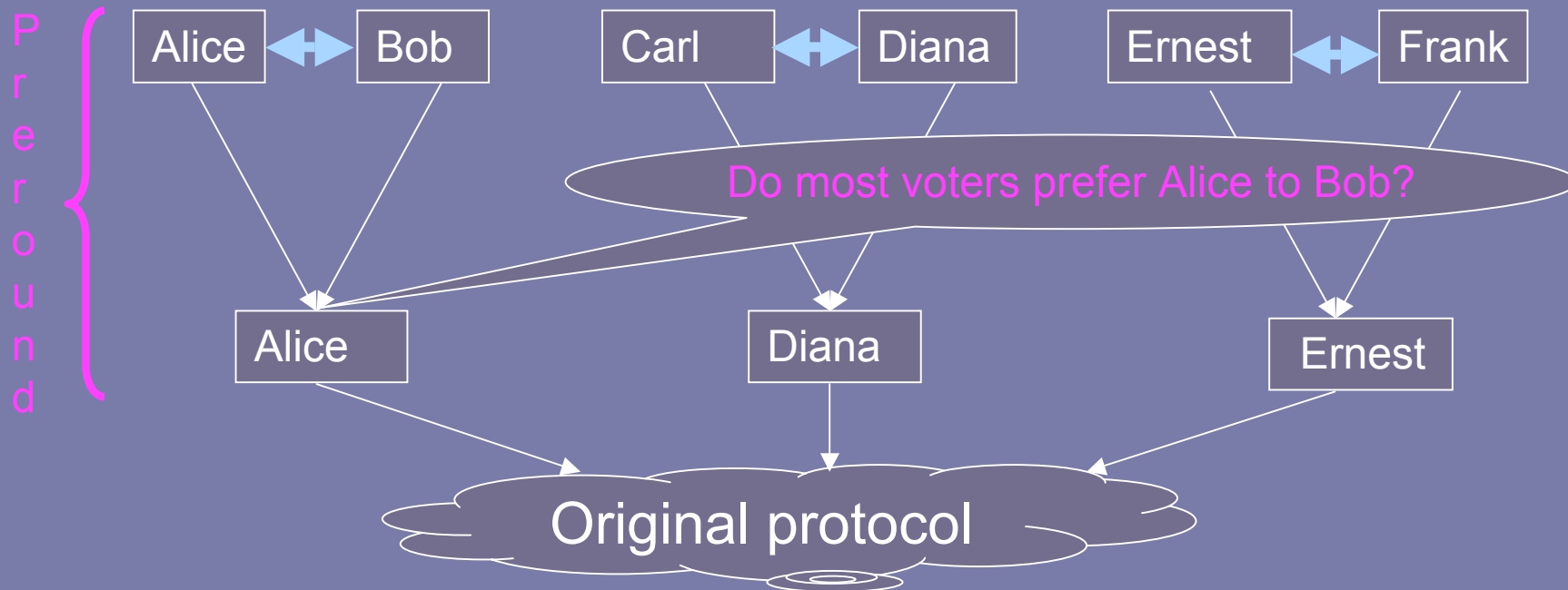
# How Do We Get Around The Impossibility Result?

- We cannot make manipulation impossible...
- But we can try to make it hard!
- How do you manipulate **Plurality**?
  - vote for your favorite candidate among those tied for the top position.
- How do you manipulate **Borda**?
  - rank your favorite **feasible** candidate highest, move his competitors to the bottom of your vote.
- How do you manipulate **STV**?
  - try all **m!** possible ballots...

# What Is Known?

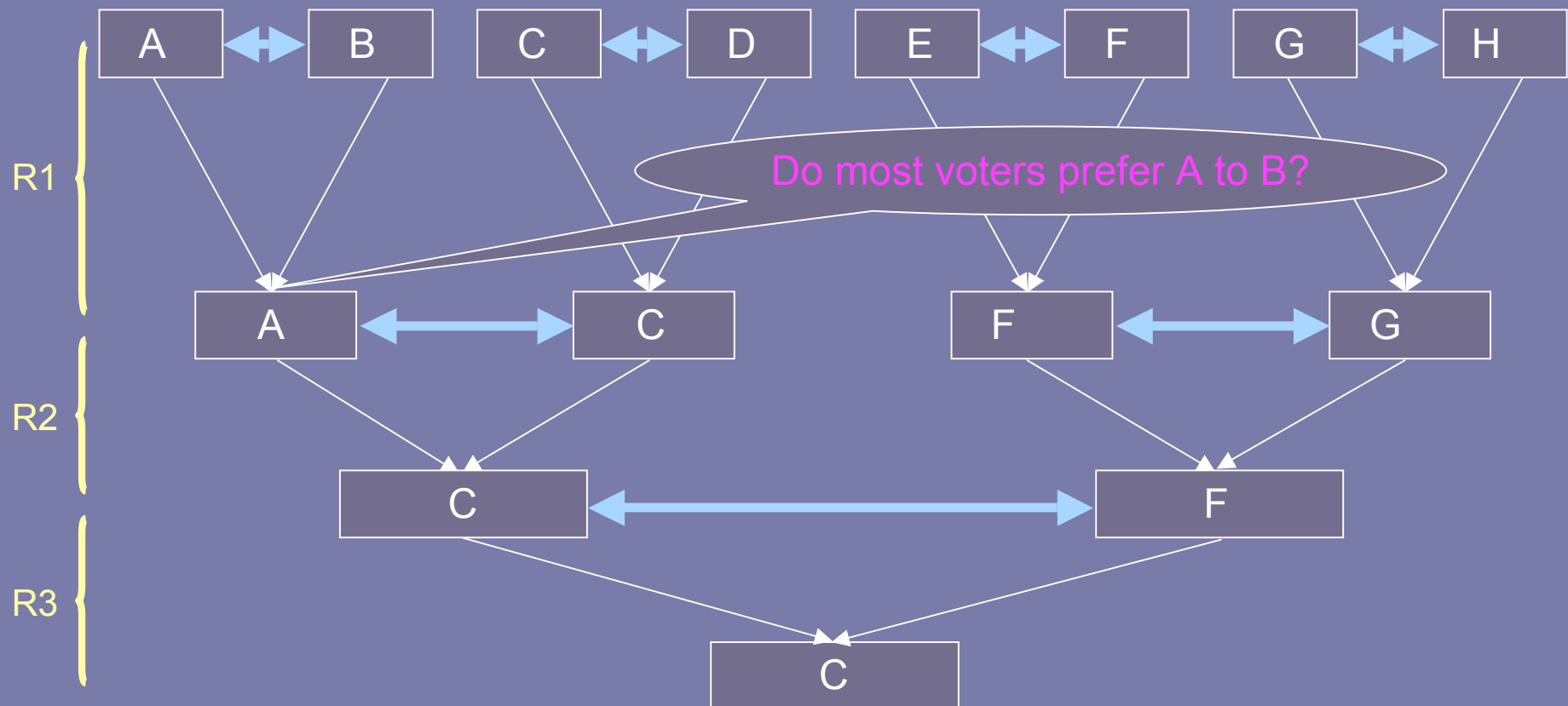
- 2<sup>nd</sup> order Copeland is NP-hard to manipulate (Bartholdi, Tovey, Trick 1989)
- STV is NP-hard to manipulate (Bartholdi, Orlin 1991)
- These rules may not reflect the welfare goals (why so many voting rules out there?)
- Want a universal method to turn any voting protocol into a hard-to-manipulate one.

# Adding a Preround (Conitzer-Sandholm'03)



- Retains **some** of the flavor of the original protocol.
- Is **NP-hard** to manipulate for many base protocols.
- Still, the outcome may be very different from the original protocol...

# Binary Cup



Binary Cup itself is easy to manipulate.

# Our Work: Hybrid Protocols

- Protocols with a preround can be viewed as **hybrids** of **BC** and other protocols
  - how about other hybrids?
- **Hyb**( $X_k$ ,  $Y$ ): execute  $k$  steps of  $X$ , then apply  $Y$  to the remaining candidates.
  - **step**: protocol-dependent
    - round of **STV** or **BinaryCup**
    - eliminating the lowest scoring candidate for **Plurality**, **Borda**
  - **Hyb**(**Plurality** $_k$ , **Borda**):
    - eliminate  $k$  candidates with the lowest **Plurality** scores
    - compute **Borda** scores w.r.t. survivors.

# New Protocols

- $\text{Hyb}(X_k, \text{STV}), \text{Hyb}(\text{STV}_k, Y)$  are NP-hard to manipulate (for any reasonable  $X, Y$ )
  - is  $\text{Hyb}(X_k, Y)$  non-manipulable for any  $X$  (or  $Y$ ) that is non-manipulable?
- $\text{Hyb}(\text{Borda}_k, \text{Plurality})$   
is NP-hard to manipulate
- $\text{Hyb}(\text{Maximin}_k, \text{Plurality})$   
is NP-hard to manipulate

# Hybrid of a Protocol with Itself

- Generally,  $\text{Hyb}(X_k, X) \neq X$ 
  - (and may be much harder to manipulate)
- $\text{Hyb}(\text{Plurality}_k, \text{Plurality})$ :
  - eliminate  $k$  lowest-scoring candidates
  - recompute the scores
  - select  $\text{Plurality}$  winner wrt  $\text{new}$  scores
- $\text{Hyb}(\text{Plurality}_1, \dots, \text{Plurality}_m) =$
- $\text{Hyb}(\text{Borda}_k, \text{Borda})$ 
  - is  $\text{NP-hard}$  to manipulate



# Limitations and Extensions

- Is  $\text{Hyb}(X_k, Y)$  hard to manipulate for any  $X, Y$ ?
  - NO:  $\text{Hyb}(\text{Plurality}_k, Y)$  is almost as easy to manipulate as  $Y$
- Utility-based voting (voters **rate** candidates rather than **rank** them)
  - HighScore: the candidate with **max** total score wins
  - manipulating  $\text{Hyb}(\text{HighScore}_k, \text{HighScore})$  is NP-hard