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# The Discursive Dilemma as a Lottery Paradox

Igor Douven
Institute of Philosophy
University of Leuven

Jan-Willem Romeijn
Department of Psychology
University of Amsterdam

# **Outline of presentation**

- The discursive dilemma
- 2 The lottery paradox
- Somorphy of the paradoxes
- A new impossibility result

#### O Discursive dilemma

Imagine a parliament of three members, voting on the following policy statements.

| voter<br>statement     | $v_1$ | $v_2$ | $v_3$ | majority |
|------------------------|-------|-------|-------|----------|
| $A_1$                  | 1     | 0     | 1     | 1        |
| $A_2$                  | 0     | 1     | 1     | 1        |
| $\neg (A_1 \land A_2)$ | 1     | 1     | 0     | 1        |

If the collective profile is assumed to be closed under conjunction, it is inconsistent.

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# Impossibility theorem

List and Pettit [2002] prove roughly the following:

The following conditions on judgment aggregation are jointly inconsistent:

- (1) the agenda has at least two independent propositions;
- (2) voters have universal domain and anonimity;
- (3) the voting rule R satisfies independence and neutrality;
- (4) R leads to consistent and complete collective opinions.

There have been many refinements of this result, but we employ this early version as base case.

# 2 Lottery paradox

Now imagine that we are considering propositions  $A_i$  stating that ticket i will lose in a lottery of three tickets:

| statement                        | prob | accept |
|----------------------------------|------|--------|
| $A_1$                            | 2/3  | 1      |
| $A_2$                            | 2/3  | 1      |
| $A_3$                            | 2/3  | 1      |
| $\neg (A_1 \land A_2 \land A_3)$ | 1    | 1      |

Assuming that we also accept the deductive closure of accepted propositions, the rule  $Accept(\varphi)$  if  $Prob(\varphi) > \frac{1}{2}$  gives inconsistent sets of accepted propositions.

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# **Acceptance rules**

Take any value for the threshold t in the rule  $Accept(\varphi)$  if  $Prob(\varphi) > t$ . There is always a sufficiently large lottery to generate inconsistency.

Accept  $\varphi$  if  $Prob(\varphi) > t$ , unless some formally specified defeater  $D(\varphi)$  holds.

Example:  $D(\varphi)$  holds if  $\varphi$  is included in some minimal inconsistent set of  $\psi_i$  for which  $Prob(\psi_i) > t$ .

If we want to maintain that sets of accepted propositions are the deductively closed, we must an acceptance rule of the above kind, which incorporates further conditions.

#### Structural acceptance

Douven and Williamson [2006] proved a general result on the lottery paradox concerning acceptance rules with defeaters, of which we use the following corrollary.

The following conditions on rational acceptance of propositions  $\phi$  are jointly inconsistent:

- (1) the possible worlds interpreting the propositions φ are equally probable;
- (2) the acceptance rule defines a structural property;
- (3) the accepted propositions are consistent, closed under conjunction, and include  $\varphi$  with  $Prob(\varphi) < 1$ .

#### No strictly formal solution

The result on structural acceptance is quite general. It covers all rules that can be defined in (higher order) logic, set theory, etc.

A function f over propositions φ is an automorphism iff

(1) 
$$f(\phi \land \psi) = f(\phi) \land f(\psi)$$
;

(2) 
$$f(\neg \varphi) = \neg f(\varphi)$$
;

(3) 
$$Prob(\varphi) = Prob(f(\varphi))$$
.

A property A of propositions  $\phi$  is structural iff it is invariant under all automorphisms f.

This also means that excluding the acceptance of inconsistent conjunctions of accepted propositions does not help.

#### S Isomorphic paradoxes

Note that we can represent the probability assignment figuring in the lottery paradox by means of equally probable possible worlds.

| world statement                  | $\mathbf{w}_1$ | $\mathbf{w}_2$ | $W_3$ | prob | accept |
|----------------------------------|----------------|----------------|-------|------|--------|
| $A_1$                            | 0              | 1              | 1     | 2/3  | 1      |
| $A_2$                            | 1              | 0              | 1     | 2/3  | 1      |
| $A_3$                            | 1              | 1              | 0     | 2/3  | 1      |
| $\neg (A_1 \land A_2 \land A_3)$ | 1              | 1              | 1     | 1    | 1      |

A representation of a probability assignment over the propositions  $\phi$  in terms of equiprobable worlds can always be given.

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#### Worlds are voters

Possible worlds can be considered as anonimous voters. The equal probability of the worlds translates into the equal say that voters have in the collective opinion.

| voter                            | $v_1$ | $v_2$ | $v_3$ | vote | accept |
|----------------------------------|-------|-------|-------|------|--------|
| $A_1$                            | 0     | 1     | 1     | 2/3  | 1      |
| $A_2$                            | 1     | 0     | 1     | 2/3  | 1      |
| $A_3$                            | 1     | 1     | 0     | 2/3  | 1      |
| $\neg (A_1 \land A_2 \land A_3)$ | 1     | 1     | 1     | 1    | 1      |

The acceptance rule  $Accept(\varphi)$  if  $Prob(\varphi) > \frac{1}{2}$  then is a majority vote.



# **Employing the isomorphy**

We want to use the result on rational acceptance rules as an impossibility theorem concerning voting rules. For this we must establish the following translations.

- Acceptance rules  $Accept(\varphi)$  translate into aggregation rules  $R(\varphi)$ .
- Because possible worlds translate into voters, these voters are essentially characterised by their opinion profile. So voters cannot have identical profiles.
- Relatedly, the voting agenda consists of the powerset of all voters.



# Agenda and domain assumptions

Both the interplay between agenda and voters and the fact that the voters are identifiable by their profiles require some further explanation.

- As opposed to other impossibility theorems, the present result employs a fixed profile to derive the inconsistency.
- The voting body may also be divided into equal parties with identifiable profiles. Such voting bodies are called party-wise opinionated.
- The agenda consisting of the powerset of parties may still be unusually rich. On the flip side, this enables us to widen the scope of voting rules significantly.

# **4** A new impossibility result

The translation leads to the following generalised impossibility result, in which votes need not be independent.

The following conditions on voting rules are jointly inconsistent:

- (1) the agenda allows for party-wise opinionated profiles;
- (2) the domain of the voting rule consists of these profiles;
- (3) the voting rule satisfies structuralness;
- (4) the collective opinion profile is consistent, closed under conjunction, and it includes propositions that are not unanimously accepted.



#### Relations to other results

The conditions in this theorem relate in rather intricate ways to the conditions of other theorems, and this requires explicit attention.

- The theorem concerns the possibility of consistent collective opinion at specific points in the domain of the voting rule. Thus unanimity need only apply at those points.
- The impossibility result nevertheless reflects back on voting rules in general because we cannot at the onset exclude these specific points.
- The condition of structuralness entails that the voters are anonimous, and further that the voting rule is neutral with respect to types of propositions.



#### **Discussion**

We conclude with some considerations on the isomorphy and the impossibility result that can be derived from it.

- The main quality of the present result is that it allows votes on propositions to be interdependent. We can drop this assumption because we assume a rich agenda.
- We may expand voting rules to include non-formal properties of propositions, such as modal notions or semantics. This is perhaps bad for rational acceptance, but fine for voting.
- Given the liveliness of the judgment aggregation literature, there may very well be applications of the isomorphy in opposite direction.